Mind the Gap: Ceremonies for Applied Secret Sharing

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Secret Sharing: (t, n) - Threshold Schemes



Secret s Size n group Threshold t

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Properties of (t, n) - Threshold Scheme

 Reconstruction: any size t subset of the n participants can compute the secret given their t shares

Secrecy: no subset of the n participants consisting of t-1 or fewer participants is able to gain any knowledge of the secret given their combined shares



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Ceremonies and Secret Sharing





Ellison, 2007

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Our Secret Sharing Ceremony Analysis Framework

- 1. Identify the stages of the ceremony
- 2. Define the threat model
- 3. Define the mode of operation
- 4. Evaluate the security goals against the adversaries

Case Study: Sunder

- → A tool from Freedom of the Press for journalists
- → Implements Shamir secret sharing
- → Support for share integrity
- → (Some) support for Base and Extended modes

• • •	Sunder	
👽 Si	Inder Create Secret Shard	
< Back		
ENTER YO	DUR SECRET	■ Hide secret?
SHARE OF 2 Sh 2 of 3 sha	PTIONS ares needed to cover the secret all Total n shares ares needed to recover secret.	number of S
	Create Secret Share	es

Sunder Stages and Modes

- → Secret Preparation
- → Share Generation
- Share Distribution
- → Secret Reconstruction
- Extended Reconstruction





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1. Choice: Select **n** participants

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- 2. Choice: Select a secure communication channel



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3. Action: The dealer sends each participant their share and

corresponding public verification key

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4. Action: Delete each share from the dealer's device.

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- **4.** Action: Delete each share from the dealer's device.
- 5. Choice: Each participant selects an appropriate storage mechanism for their share

- **1**. Choice: Select **n** participants
- 2. Choice: Select a secure communication channel



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3. Action: The dealer sends each participant their share and

corresponding public verification key

- 4. Action: Delete each share from the dealer's device.
- 5. Choice: Each participant selects an appropriate storage mechanism for their share
- **6.** Action: Each participant stores their share

Sunder: Analysis Threat Model



- → A high-powered adversary with the power and resources of a government actor
- → Adversaries may be **participants or outsiders**
- → We do not assume roles are static
- Adversarial goals may include: learning secret information, modifying secret information, preventing secret recovery, and causing harm to participants

Sunder Ceremony Evaluation

●=achieved; ①=ceremony dependent; ○=not achieved

	Classic Shamir				Sunder Ceremony			
	Base		Ext		Base		Ext	
	HBC	MAL	HBC	MAL	HBC	MAL	HBC	MAL
t-Sep. Priv.	•	•	•	•	•	•	•	•
Availability	•	•	0	0	•	•	O	O
IT Sec.	O	Ð	0	0	0	0	0	0
Conf.	O	Ð	Ð	O		Ð	Ð	D
Integrity	0	0	0	0	•	•		O

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Threats to Secret Reconstruction

- 1. Alice leaving the organization
- 2. A share being damaged
- 3. A share being stolen
- The device storing the encrypted files is destroyed



Our Lightweight Proactive VSS

Adds three new stages:

- → Share Update
- Share Validate
- → Generate Commitment



Access Revocation via Updates



Verification of Share Integrity and File Integrity

Proactive VSS: Share Validation

- Action: The participant fetches the commitment from its trusted public location
- 2. Device: The participant will evaluate the validation function
- 3. Device: The participant verifies the correctness of her share by checking the commitment matches the validation function



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Lightweight Improvements Comparison

●=achieved; ①=ceremony dependent; ○=not achieved

	Classic Shamir				Our Proactive VSS			
	Base		Ext		Base		Ext	
	HBC	MAL	HBC	MAL	HBC	MAL	HBC	MAL
t-Sep. Priv.	•	•	•	•	•	•	•	•
Availability	•	•	0	0	•	•	•	•
IT Sec.	O	Ð	0	0	0	0	0	0
Conf.	O	Ð	Ð	O	•	•	•	•
Integrity	0	0	0	0	•	•	•	•

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Takeaways

- → We present a framework to facilitate the analysis of practical threshold schemes
- → Variations in the ceremony can lead to changes in the fundamental security properties provided to end users
- → Our framework can aid in the design and analysis of future implementations of secret sharing through its detailed ceremony definition and explicit coverage of previously undefined assumptions

Sunder Stages: Share Generation

- 1. Choice: The dealer chooses values for **n** and **t**
 - 2. Device: Generates a signature key pair
 - 3. Device: Generates n shares
- 4. Action: Delete all copies of **s** and the signature key. (The device retains the public verification key.)

Extended Mode



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Shatter Secrets: Evaluation

●=achieved; ①=ceremony dependent; ○=not achieved

		Classic	Shatter Secrets			
	Base		Ext		Ext	
	HBC	MAL	HBC	MAL	HBC	MAL
t-Sep. Priv.	•	•	•	•	•	٠
Availability	•	•	\bigcirc	0	0	\bigcirc
IT Sec.	0	O	0	0	0	0
Conf.	0	O	Ð	O		
Integrity	0	0	0	0		O

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Proactive VSS: Share Validation

1. Action: The participant fetches the commitment \dot{C} from its trusted public location **2**. Device: Using ϕ_0,\ldots,ϕ_{t-1} which constitute \dot{C} , the participant will then calculate ψ by evaluating $\prod_{i=0}^{t-1} \phi_i^{{a_i}^j}$ Device: The participant can then validate the correctness of her

share by validating $\,\psi$ is equal to $\,g^{f(a_i)}\,$